
Nuprl Tutorial

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Outline

- What are:
 - the goals
 - the philosophy
 - the domain
 - *of NuPRL*
- Interactive Theorem Proving
- System
 - Architecture
 - Logic “engine”
 - User interface

The screenshot shows a theorem prover interface with a central proof window and a list of theorems on the right. The central window displays the following proof:

```

- top
+ F  $\forall i, j, l: \mathbb{Z}. \forall k: \{l: \mathbb{Z} \mid l \neq 0 \in \mathbb{Z}\}.$ 
   $i * k = j * k + l \text{ rem } k \in \mathbb{Z} \Rightarrow i - j = 0 \in \mathbb{Z}$ 
BY UnivCD
  THENW Auto
1- 1.  $i: \mathbb{Z}$ 
   2.  $j: \mathbb{Z}$ 
   3.  $l: \mathbb{Z}$ 
   4.  $k: \{l: \mathbb{Z} \mid l \neq 0 \in \mathbb{Z}\}$ 
   5.  $i * k = j * k + l \text{ rem } k \in \mathbb{Z}$ 
  F  $i - j = 0 \in \mathbb{Z}$ 
    
```

The right-hand pane lists several theorems, including:

- *t rem_bnd_thm3 $\forall a: \{i: \mathbb{Z} \mid (i < 0)\}. \forall b$
- *t div_rem_sum_thm $\forall i: \mathbb{Z}. \forall j: \{k: \mathbb{Z} \mid k \neq 0$
- T divides_thm1 $\forall i, j, l: \mathbb{Z}. \forall k: \{l: \mathbb{Z} \mid l$
- *t divides_thm2 $\forall i, j, l: \mathbb{Z}. \forall k: \{l: \mathbb{Z} \mid l$
- *t div_zero_thm $\forall i: \{i: \mathbb{Z} \mid i \neq 0 \in \mathbb{Z}\}$
- *t int_div_mul_thm $\forall i: \mathbb{Z}. \forall j: \{k: \mathbb{Z} \mid k \neq 0$
- *t int_mul_zero_thm $\forall i, j: \mathbb{Z}. i * j = 0 \in$
- *t int_lt_thm $\forall n, m: \mathbb{Z}. 0 < n - m \Rightarrow$
- *t int_gt_thm $\forall n, m: \mathbb{Z}. n - m < 0 \Rightarrow$
- *t div_ident $\forall a: \mathbb{Z}. a \div 1 = a \in \mathbb{Z}$
- *D power_df Parends :: Prec(div)::
- *M power_m1 % Take n to the powr
- *t power_wf $\forall n: \mathbb{N}. \forall m: \mathbb{Z}. n^m \in \mathbb{N}$
- *t comb_for_power_wf $(\lambda n, m, z. n^m) \in n: \mathbb{N} \rightarrow$
- *t power_wf2 $\forall n, m: \mathbb{Z}. n^m \in \mathbb{Z}$

Philosophy

- Formalize & implement *mathematics* and *computation*
- *Assist* the implementor
- Logic/type theory is the formal language of choice
 - Possible to derive algorithms without ever writing a program
 - Programming comes as a by-product
- Truth of statements is undecidable
 - Proofs are by interaction
 - A great deal of *assistance* is available (checking, prompting, documenting)
 - Some *automation* available
 - System goals:
 - ◆ develop automation as well as math domains
 - ◆ make the user interface comfortable

Logic & Type Theory

- How can mathematics be formalized?
- Use a logic with statements, assertions, and inference rules
- Constructive type theory: higher order logic + computation
 - Types provide specifications for programs
 - Types are assertions
 - Programs can be derived from proofs, or can be shown to inhabit types
- Types (specifications/assertions/propositions):
 - Void, Int, * list
 - Equality
 - Functions ($x:A \rightarrow B$), Products ($x:A \times B$), Disjoint Unions ($A + B$)
 - Type universes (U_i is all types at level i , $U_i \in U_{i+1}$)

Types I

- Function space
 - $\forall i:\mathbb{N}. i \geq 0$
 - $i:\mathbb{N} \rightarrow i \geq 0$
 - $i:\mathbb{N} \rightarrow \{i \dots\}$
- Product space
 - $\exists i:\mathbb{N}. \text{TM}(i) \text{ halts}$
 - $\forall i:\mathbb{N}. \exists j:\mathbb{N}. j * j \leq i \wedge (j + 1) * (j + 1) > i$
 - $i:\mathbb{N} \rightarrow j:\mathbb{N} \times (j * j \leq i) \times ((j + 1) * (j + 1) > i)$
- Disjoint union
 - $A + B$
 - $\forall i:\mathbb{N}. \text{TM}(i) \text{ halts} \vee \neg \text{TM}(i) \text{ halts}$

Types II

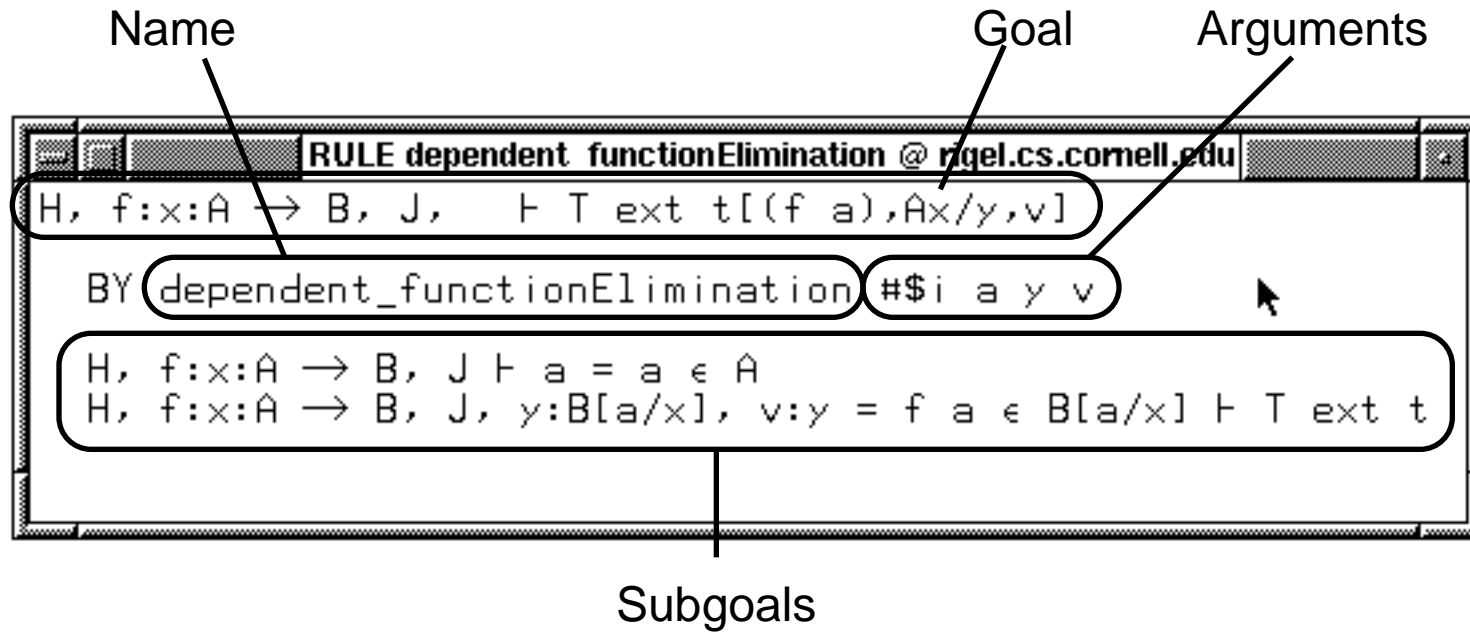
- Type universes
 - U_i is all types/propositions to level i
 - $Z, \text{Void} \in U_1$
 - $A, B \in U_i \Rightarrow x:A \rightarrow B \in U_i$, etc
 - $\forall P:N \rightarrow U_i. P(0) \Rightarrow (\forall i:N. P(i) \Rightarrow P(i+1)) \Rightarrow (\forall i:N. P(i))$
 - $U_i \in U_{i+1}$
- Set
 - $\text{car}:U_i \times \text{insert}:Z \rightarrow \text{car} \times \text{member}: Z \rightarrow \text{car} \rightarrow U_i \times \dots \in U_{i+1}$
- Equality types
 - $a = b \in T$
 - T is a well-formed type
 - a and b are well-formed elements of T
 - a and b are equal elements in T
 - membership: $a = a \in T$ means $a \in T$

Sequents

- Sequent: $a_1:H_1, a_2:H_2, \dots, a_n:H_n \mid\!\!\!-\ G$
 - H_1 is a type
 - for any $a_1 \in H_1$, $H_2[a_1]$ is a type
 - ...
 - for any $a_1 \in H_1, \dots, a_{n-1} \in H_{n-1}[a_1, \dots, a_{n-2}]$, $H_n[a_1, \dots, a_{n-1}]$ is a type
 - for any $a_1 \in H_1, \dots, a_n \in H_n[a_1, \dots, a_{n-1}]$, $G[a_1, \dots, a_n]$ is a type, and it is true
- Functionality
 - for any $a_1, b_1 \in H_1$, $H_2[a_1]$ and $H_2[b_1]$ are equal types
 - ...

Rules

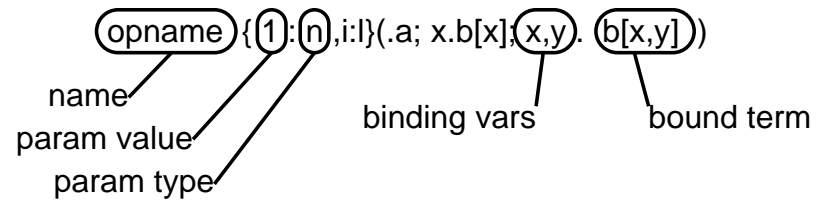
- A rule is an implication on sequents:



Syntax

- Everything is a *term*:

- $\text{term} ::= \text{opname}\{\text{params}\}(\text{bterms})$
- $\text{opname} ::= \langle \text{string} \rangle$
- $\text{params} ::= \epsilon \mid \text{paramlist}$
- $\text{paramlist} ::= \text{param} \mid \text{paramlist} , \text{param}$
- $\text{param} ::= \langle \text{number} \rangle : n \mid \langle \text{string} \rangle : s \mid \text{level-exp} : l \mid \dots$
- $\text{bterms} ::= \epsilon \mid \text{btermlist}$
- $\text{btermlist} ::= \text{bterm} \mid \text{btermlist} ; \text{bterm}$
- $\text{bterm} ::= \text{vars} . \text{term}$
- $\text{vars} ::= \epsilon \mid \text{varlist}$
- $\text{varlist} ::= \text{var} \mid \text{varlist} , \text{var}$
- $\text{var} ::= \langle \text{string} \rangle$



Syntax Examples

- Number: $1 \equiv \text{natural_number}\{1:n\}$
- Variable: $x \equiv \text{variable}\{x:v\}$
- Summation: $i + j \equiv \text{add}\{ \}(\text{.variable}\{i:v\}; \text{.variable}\{j:v\})$
- Abstraction: $\lambda x.x + 1 \equiv \text{lambda}\{ \}(x.x + 1)$

User Model

- The user builds domains by:
 - creating definitions of:
 - ◆ types
 - ◆ methods
 - ◆ properties
 - verifying properties in *theorems*
- The interface is interactive
 - As opposed to batch AI first order theorem provers
 - Premises are complex
 - Backward chaining (goal oriented)
 - ◆ Supply the goal of a theorem
 - ◆ *Refine* it with a rule or a tactic to generate subgoals
 - ◆ Prove the subgoals

Example Proof Step

```

THM stamps_thm @ rigel.cs.cornell.edu
# top 1 2 1
1. postage : ℤ
2. 0 < postage
3. postage - 1 > 8
   ⇒ (∃i,j:ℕ. 3 * i + 5 * j = postage - 1 ∈ ℤ)
4. postage ≥ 8
   ⊢ ∃i,j:ℕ. 3 * i + 5 * j = postage ∈ ℤ

BY Decide 「postage = 8 ∈ ℤ」
   THENW Auto

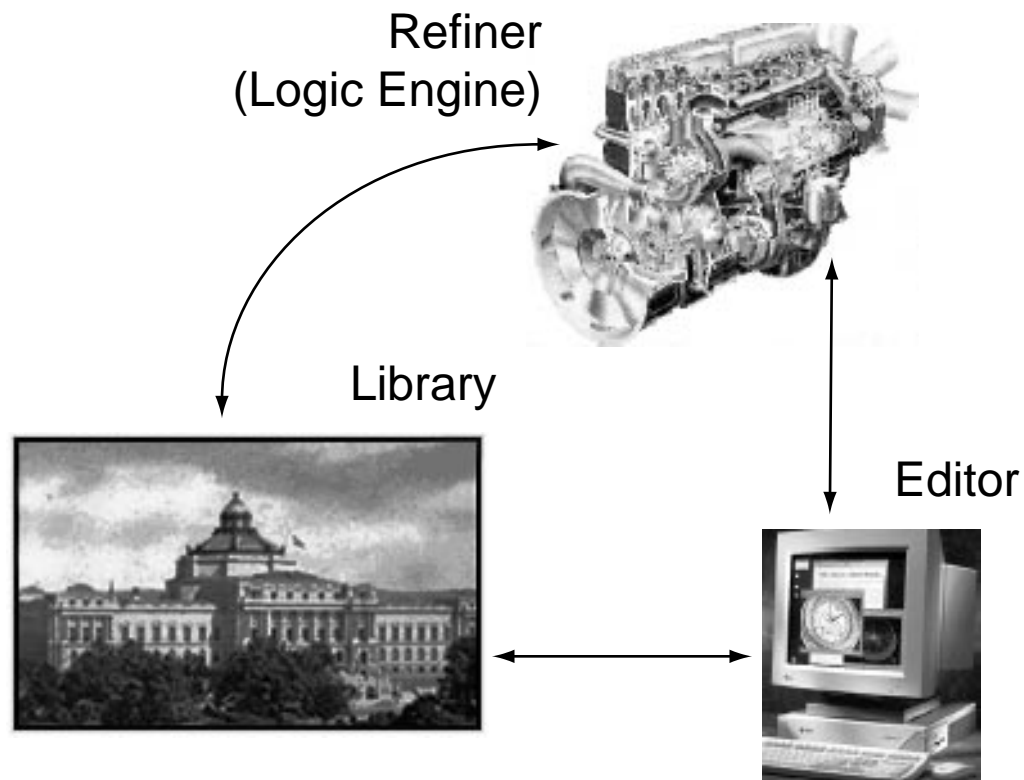
1# 5. postage = 8 ∈ ℤ
   ⊢ ∃i,j:ℕ. 3 * i + 5 * j = postage ∈ ℤ

2# 5. ¬(postage = 8 ∈ ℤ)
   ⊢ ∃i,j:ℕ. 3 * i + 5 * j = postage ∈ ℤ

```

System Architecture

- Refiner enforces the logic
- Library maintains database of definitions and inference steps
- Editor provides a user interface



Refiner

- Performs refinement according to rule collection



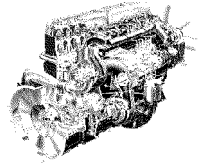
```
THM excluded_middle_thm @ rigel.cs.cornell.edu
# top 1
1. P :  $\mathbb{P}$ 
2.  $\neg\neg P$ 
 $\vdash P + \neg P$ 

BY Refine `inrFormation`
    [mk_level_exp_arg (mk_level_exp ['i', 0])]

1#  $\vdash \neg P$ 
2#  $\vdash P = P \in \mathbb{U}$ 
```

Tactics

- Refiner may also provide a *tactic* language
- In NuPRL, the language is ML
- Tactics can be programmed to analyze goals, and perform more intelligent tasks



```
THM excluded_middle_thm @ rigel.cs.cornell.edu
# top 1
1. P : P
2. ¬¬¬P
┆ P + ¬P

BY Sel 2 (D 0)
    THENW Auto

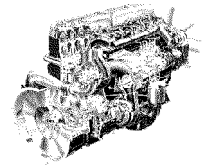
1# ┆ ¬P
```

Styles of reasoning in NuPRL

- Proofs are interactive, and refined
- Prove something is true (inhabited by some program)
- Prove something is well-formed (that it is a sensible assertion)
 - Well-formedness is a major component of proving
 - The well-formedness checking is delayed
 - A theorem must be verified to be well-formed as it is proved
 - *Example:* to prove $A \Rightarrow B$, prove A is a proposition, then assume A and prove B
 - This style differs from other major type theories
 - Well-formedness is undecidable, but the type system is quite expressive
 - Not true that if a term is well-formed, then so is every subterm
 - ◆ $\text{Void} \rightarrow (1 + \mathbb{Z})$

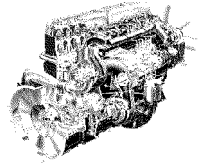
Provided Tactics

- `Auto : tactic`
 - Performs default well-formedness reasoning (gets most cases)
 - Performs simple logical reasoning
 - Tries to limit search so that every step gets closer to a proof
- `D : int → tactic`
 - “Decompose” a clause
 - on 0, performs an *introduction*
 - on a hyp, performs an *elimination*, case analysis or induction
- `NatInd : int → tactic`
 - better form of elimination on natural numbers
- `Arith : tactic` and `SupInf : tactic`
 - Perform reasoning about arithmetic
 - Linear systems of equations
- `RW : conv → int → tactic`
 - Rewrites using equivalences or implications



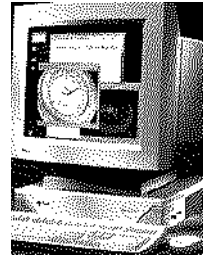
Provided Tacticals

- A tactical is a function taking a tactic as an argument
- Tac1 **THEN** Tac2
 - Run Tac1 on goal, then run Tac2 on all subgoals
- Tac1 **THENW** Tac2
 - Run Tac1 on goal, then run Tac2 on all well-formedness subgoals
- Tac1 **ORELSE** Tac2
 - Run Tac1 on goal. If it fails, run Tac2
- Also have:
 - functions to examine terms
 - functions to build terms
 - all standard functions in ML
- Tacticals are secure because there is no way to construct a tactic, except from other tacticals.



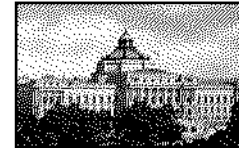
Editor

- What are the objects of the system?
 - Definitions
 - Rules
 - Theorems
 - Comments
 - Code
- All of them contain terms
- Structured editing of terms
- Each term may have a *display form*



Library

- Library window displays loaded objects
- Organized into *theories*
 - Linear list of items
 - Begins with comment object called **name**_begin
 - Ends with comment object called **name**_end
- One liner for each item
 - **Status** *: correct and complete, #: partial, -: incorrect
 - **Type**
 - ◆ C: comment, D: display form, T: theorem, A: abstraction, R: rule, M: ML code
 - ◆ lower case if object has not been checked
 - **Synopsis** first line of contents



Library Window

```

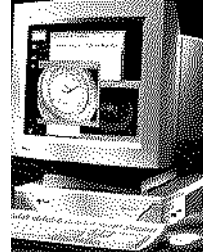
Library (File: ~/nuprl/theories/horus/event_1.thy) @ rigel.cs.cornell.edu
*C udp_eid_start      ----- begin udp_eid -----
*D udp_eid_type_df   UdpEidType == udp_eid_type{}
*A udp_eid_type      UdpEidType == addr:INetAddr × port:Port × Mux
*t udp_eid_type_wf   UdpEidType ∈  $\mathbb{U}_1$ 
*D udp_addr_df      <t:udp:E>.addr == udp_addr{}(<t>)
*A udp_addr         t.addr == t.1
*t udp_addr_wf      ∀t:UdpEidType. t.addr ∈ INetAddr
|*D udp_port_df     <t:udp:E>.port == udp_port{}(<t>)
*A udp_port         t.port == t.2.1
*t udp_port_wf      ∀t:UdpEidType. t.port ∈ Port
*D udp_mux_df       <t:udp:E>.mux == udp_mux{}(<t>)

ML top loop @ rigel.cs.cornell.edu
M> jump "state_1_begin" | ;;

```

Display Forms

- Terms may be:
 - **primitive**, like $\lambda v.b$ (lambda {} (v.b)), or $Z(\text{int}\{\}\ ())$
 - **defined**, like $\text{let } v = e \text{ in } b$
 $\text{let}\{\}(.e; v.b) \rightarrow (\lambda v.b)(e)$
- A display form may be defined for any term
- Specifies how the term should be printed



```

DISP xlet_df @ rigel.cs.cornell.edu
let <v:var> = <e:expr:*> in <b:body:*>|
== xlet{}(<e>; <v>.<b>)
  
```

- LHS contains printing directives
- Slots describe areas for terms
- When the term is constructed, the user is prompted for input in the slots
- Special instructions for line breaking, parenthesization, etc.

Abstractions

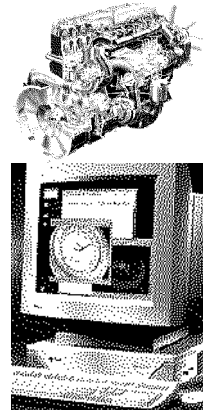
- Definitions are made through *abstractions*
- An abstraction defines a pair of terms that are:
 - the defined/definition; the redex/contractum

```
ABS xlet @ rigel.cs.cornell.edu
let v = e in b == (λv.b) e
```

- Often used to define more complex types:

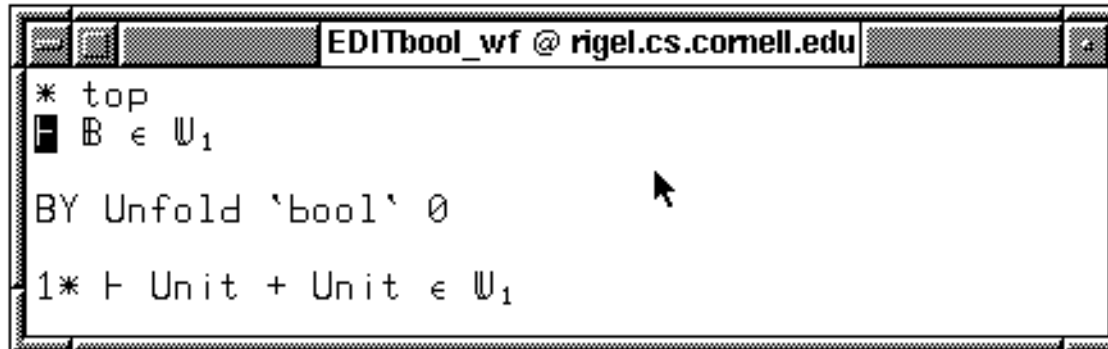
```
ABS bool @ rigel.cs.cornell.edu
B == Unit + Unit
```

```
ABS ifthenelse @ rigel.cs.cornell.edu
if b then t else f fi
== case b of inl() => t | inr() => f
```



Abstraction Well-formedness

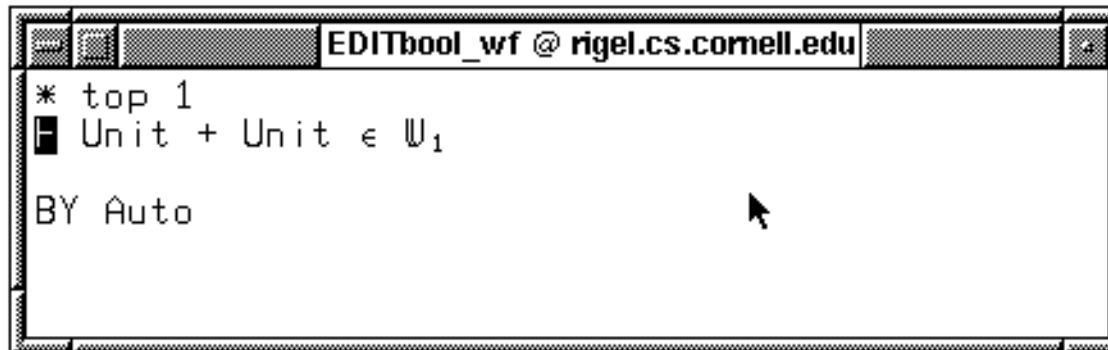
- Abstractions usually have a well-formedness theorem



```
* top
B ∈ W1

BY Unfold 'bool' 0

1* ⊢ Unit + Unit ∈ W1
```



```
* top 1
Unit + Unit ∈ W1

BY Auto
```


Larger Example

```
stab_up_hdr(event, msg, header)
== Case(<event, header>)
  Case <up_event(UpView, my_rank = my_rank', nmembers = nmembers'), NoHdr> =>
    my_rank ← my_rank';
    nmembers ← nmembers';
    acks ← make_vect2(nmembers, nmembers, 0);
    stable ← make_vect(nmembers, 0);
    ncasts ← make_vect(nmembers, 0);
    flushing ← ff;
    queue.append(stab_up(<event, msg>))
  Case <up_event(UpCast, origin = origin), NoHdr> =>
    /* Update #casts entry and set ack field */
    ncasts[origin]++;
    queue.append(stab_up(<UpSet(StabilityName,
                              [upAck(Ack(<origin, ncasts[origin]>))],
                              event), msg>))
  Case <., NoHdr> =>
    queue.append(stab_up(<event, msg>))
  Default => /* We need a way to fail */ skip
```

Large Well-formedness

- Well-formedness is quantified by types of arguments



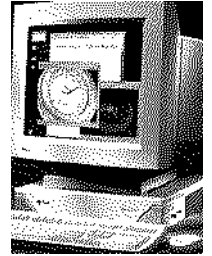
```
THM stab_up_hdr_wf @ rigel.cs.cornell.edu
# top
┆ ∀event:StabEvent. ∀msg:EventMessage.
  ∀hdr:StabHeaderType. ∀state:StabilityState.
    stab_up_hdr(event, msg, hdr) StabilityState
    ∈ StabilityState

BY UnivCD

1# 1. event : StabEvent
   2. msg : EventMessage
   3. hdr : StabHeaderType
   4. state : StabilityState
┆ stab_up_hdr(event, msg, hdr)
  StabilityState ∈ StabilityState
```

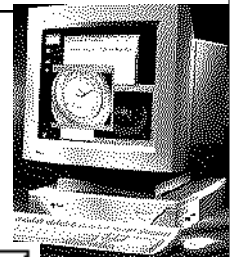
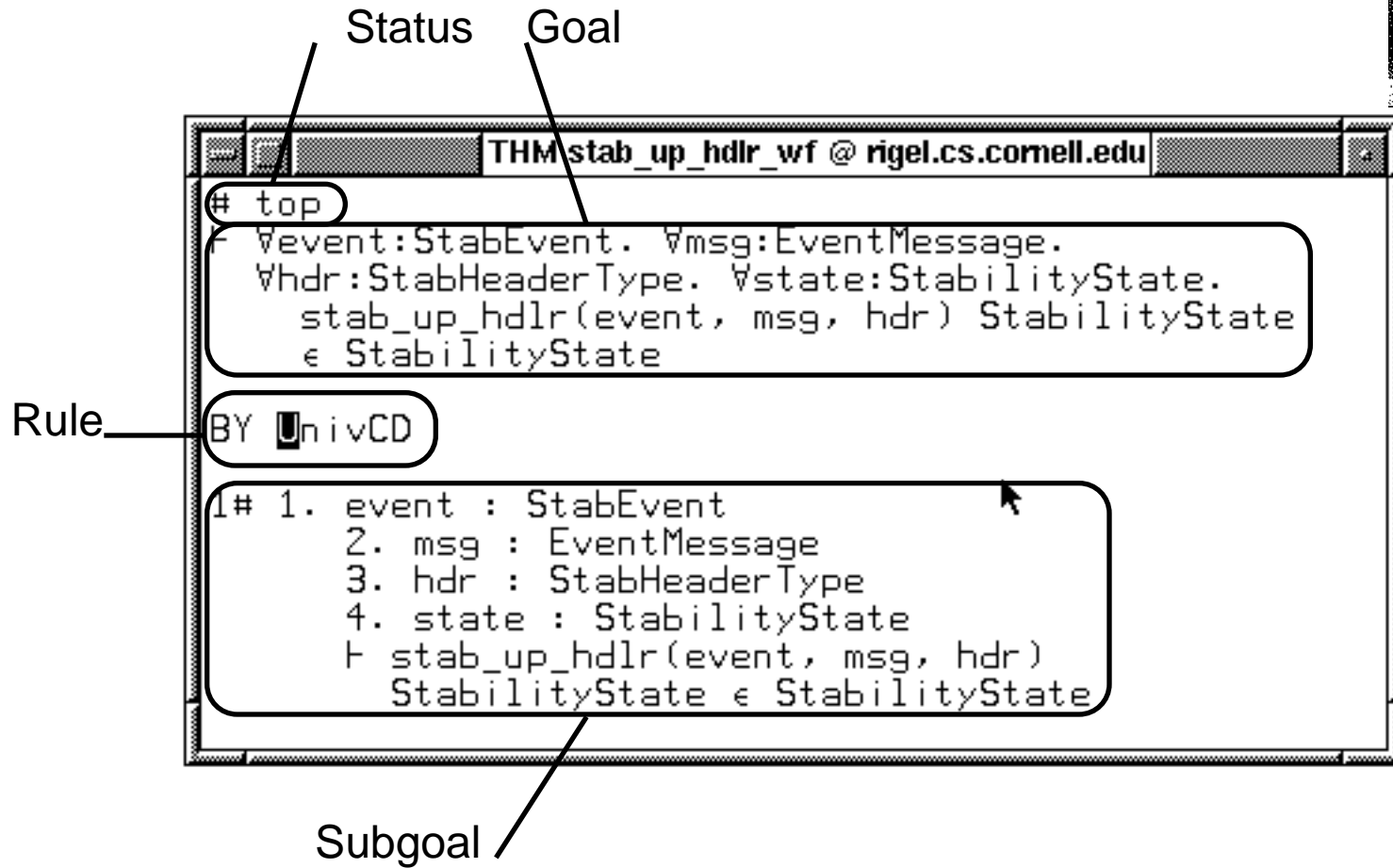
Term Editor Commands

- Large assortment of commands for editing terms
- Use either keyboard or mouse
- To get a new term, type the name of its display form (usually the operator name)
 - ^F moves *forward*, ^B moves *back*
 - ^P moves *up* the term tree
 - ^O *opens* a term slot
 - ◆ in the middle to text (like tactic text)
 - ◆ in the middle of a term (to add a subterm)
 - ^Xex *explodes* a term (displays it in canonical form)
 - ^Xim *implodes* it
 - ^K *kills* a term, ^Y *yanks* it back from the kill ring



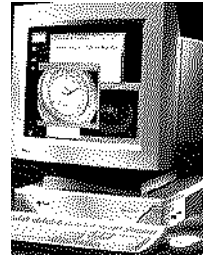
Refinement (Proof) Editor

- A window into a proof of a theorem has four parts:



Refinement

- ^O opens goal or subgoal term or rule box
- ^O edits the main goal of a theorem
- Type tactic into rule box
- ^Z closes and *checks* a term or rule box
 - Set the goal
 - Runs the tactic to produce subgoals



Summary

- Nuprl is a system for developing formal mathematics
- Formalism is based on *type theory* (logic + computation)
- System is designed to assist and automate the formalization
- Three parts:
 - Refiner (for checking and inferring proofs)
 - Library (for storing rules, definitions, and proofs)
 - Editor (human interface for editing object in the library)
- Refiner uses ML as a *tactic* language
- Editor works directly on the term tree
 - Structured editing
 - Display forms for terms

Where to go from here

- Documentation
 - Nuprl book
 - Web site (<http://www.cs.cornell.edu/Info/Projects/NuPr1/nuprl.html>)
 - ◆ Nuprl book
 - ◆ Nuprl 4.2 reference manual
 - ◆ Nuprl 4.2 tutorial
 - ◆ Library browser
- To use Nuprl
 - Execute `~nuprl/bin/run-nuprl`
 - Use large SunOS machine, like **gemini** or **virgo**
 - `xhost` the machine you are running from
 - It does run on Solaris, just not officially supported

